

SMTpp: preprocessors and analyzers for SMT-LIB

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July 18th, 2015

Introduction

SMTpp = SMT pre-processor

What? source-to-source transformer and analyzer for the SMT-LIB

- Why?**
- ▶ SMT developer: factor simplifications found in solvers $\approx 20\%$ in veriT
 - ▶ SMT developer: fuzzing-based tests, delta-debugging
 - ▶ SMT user: factor SMT configurations from generators of verification conditions
 - ▶ SMT-LIB: analyse and classify benchmarks
 - ▶ SMT-COMP: scramble benchmarks

How? high-level programming language, symbolic programming (OCaml+menhir)

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*Prototype release available from
<http://www.verit-solver.org>, under ISC licence.*

Positioning SMTpp

- ▶ SatELite: a pre-processor for SAT
 - ▶ minimization of CNF
 - ▶ target a class of SAT encodings
- ▶ jSMTLIB: SMT-LIB for Java
 - ▶ Java-based
 - ▶ SMT-LIB parser and type-checker
 - ▶ interaction with SMT-LIB and SMT-solvers
 - ▶ Eclipse plug-in
 - ▶ API
- ▶ SMT solvers (\supset SMTpp)
 - ▶ parsing
 - ▶ type inference
 - ▶ logic detection
 - ▶ simplification
- ▶ delta debugger, scrambler, fuzzer

AST

- ▶ one-to-one mapping from grammar to OCaml data type
- ▶ excerpt (commands)

```
type command_desc =  
  | CmdAssert of term  
  | CmdCheckSat  
  | CmdDeclareConst of symbol * sort  
  (* ... *)  
  
type command = {  
  command_desc : command_desc;  
  command_loc   : Locations.t;  
}
```

- ▶ Full code available at:
<https://github.com/ossanha/smtpp>

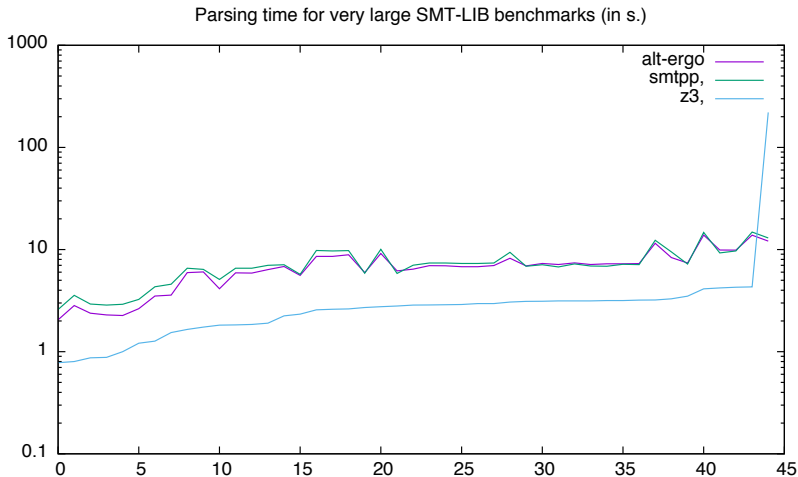
The parser module

- ▶ ocaml yacc or **Menhir**
 - ▶ LR(1) parser generator
 - ▶ keyword references to semantic values
 - ▶ improved debugging support

```
%inline command:  
  | ASSERT t=term;  
  { let loc = mk_loc $startpos $endpos in  
    mk_command (CmdAssert t) loc }  
  | CHECKSAT  
  { let loc = mk_loc $startpos $endpos in  
    mk_command CmdCheckSat loc }  
  | DECLARECONST sym=symbol; so=sort;  
  { let l = mk_loc $startpos $endpos in  
    mk_command (CmdDeclareConst(sym, so)) l }
```

```
let mk_loc loc_start loc_end =  
  { loc_start; loc_end; } ;;
```

Parser: experiments



Obfuscator module

- ▶ replace declared symbols with tool-generated names
- ▶ example input (excerpt):

```
(declare-fun x () Real)
(declare-fun f (Real) Real)
(assert (=> (> x 0) (< (f x) 0)))
```

- ▶ corresponding output:

```
(declare-fun S2 () Real)
(declare-fun S3 (Real) Real)
(assert (=> (> S2 0) (< (S3 S2) 0)))
```


Multi-script generation module

- ▶ Slice incremental benchmarks into non-incremental ones.

```
(set-logic QF_LIA)
(declare-fun w () Int)
(declare-fun x () Int)
(declare-fun y () Int)
(declare-fun z () Int)
(assert (> x y))
(assert (> y z))
(push 1)
(assert (> z x))
(check-sat)
(pop 1)
(push 1)
(check-sat)
(exit)
```

Multi-script generation module

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```
(set-logic QF_LIA)
(declare-fun w () Int)
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(declare-fun z () Int)
(assert (> x y))
(assert (> y z))
(check-sat)
(exit)
```

Use-def analysis module

- ▶ detects:
 1. use of undefined symbols: error
 2. definition of never-used symbols: noise
 - ▶ commonplace in SMT-LIB
- ▶ results verified with grep
- ▶ application
 - ▶ simplify benchmarks
 - ▶ possible errors

Use-def analysis: experiments

- ▶ BWare project: prelude with symbols not used in all files
- ▶ a number of scripts have local bindings (`let`) for variables that are not used
- ▶ all benchmarks in LRA/keymaera have a homonymous global constant symbol and existentially quantified variable.
- ▶ QF_ABV has scripts with significant amounts (30-50% file size) of unused declarations.

Logic detection module

- ▶ Errors occur when logic is not precise enough
- ▶ Logic is too restrictive:
 - ▶ reduce choice of solvers
 - ▶ might reduce performance
 - ▶ might reduce precision of results (e.g. `unknown` vs `sat`)
- ▶ Benchmark generators often choose a default logic that is sometimes too expressive
- ▶ Approaches in existing solvers
 - ▶ have a default catch-all logic
 - ▶ inspection of AST to infer a logic

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 - ▶ inspection of AST to infer a logic *
- ▶ \implies logic detector module in SMTpp.
 - ▶ Still conservative: e.g. LA benchmarks classified as NLA
 - ▶ DL: problems with non-canonical atoms, e.g. $x \leq y + 1$

Logic detection: experiments

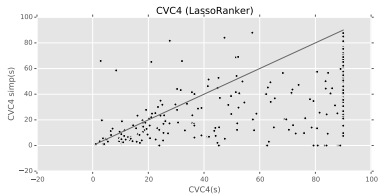
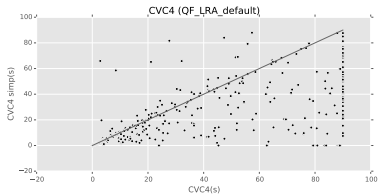
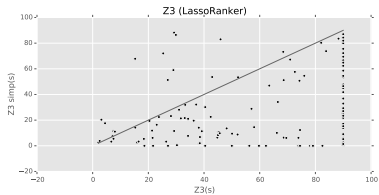
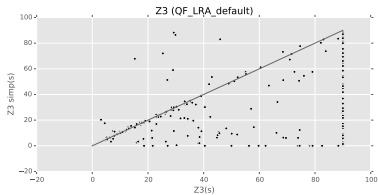
- ▶ SMT-LIB (2014 version): 5% reported by SMTpp
 - AUFLIRA 1368 scripts without Array.
 - UF/grasshopper propositional
 - QF_UFLRA 20% do not use arithmetics
 - QF_AUFLIA mostly do not use arithmetic functions, 1/4 have no free sorts and uninterpreted functions.

LA simplifications module

- ▶ motivated on LassoRanker subfamily
- ▶ straightforward transformations
 - ▶ absorbing element
 - ▶ multiplication
 - ▶ logical and
 - ▶ logical or
 - ▶ neutral element
 - ▶ multiplication
 - ▶ addition
 - ▶ turn relational constraints into boolean constants (i.e. $1 > 0$)
- ▶ Local Gauss elimination
 - ▶ identify equational constraints
 - ▶ split additive expression
 - ▶ replacing monoms
 - ▶ monoms to be replaced
 - ▶ perform substitution on relational constraints

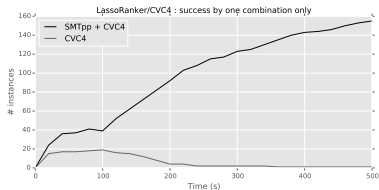
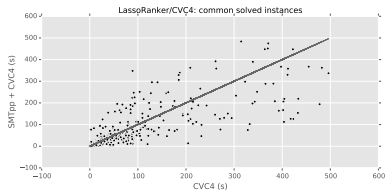
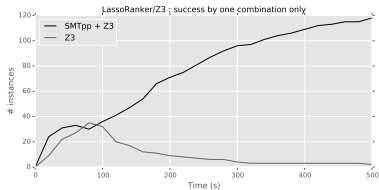
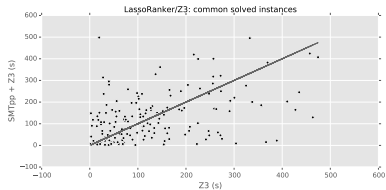
LA simplifications: experiments

Comparison of execution times



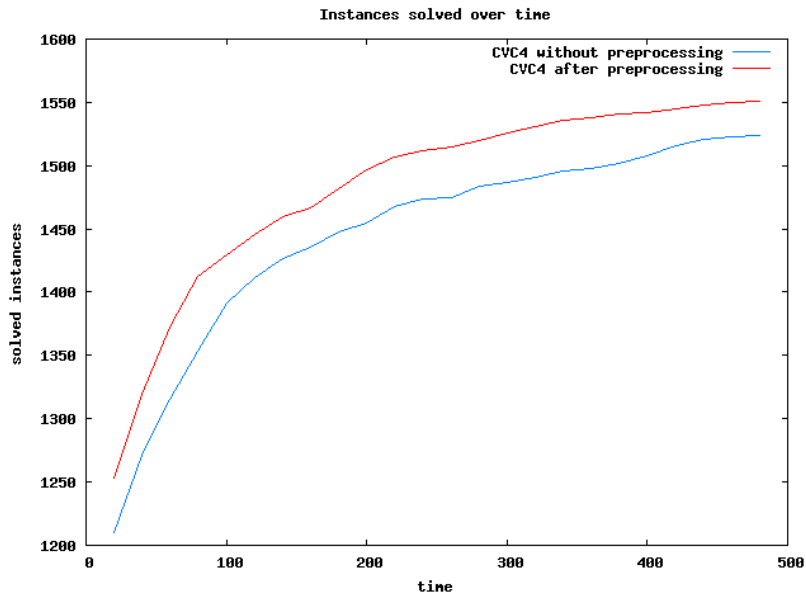
LA simplifications: experiments

Execution time on solved instances and count of uniquely solved instances



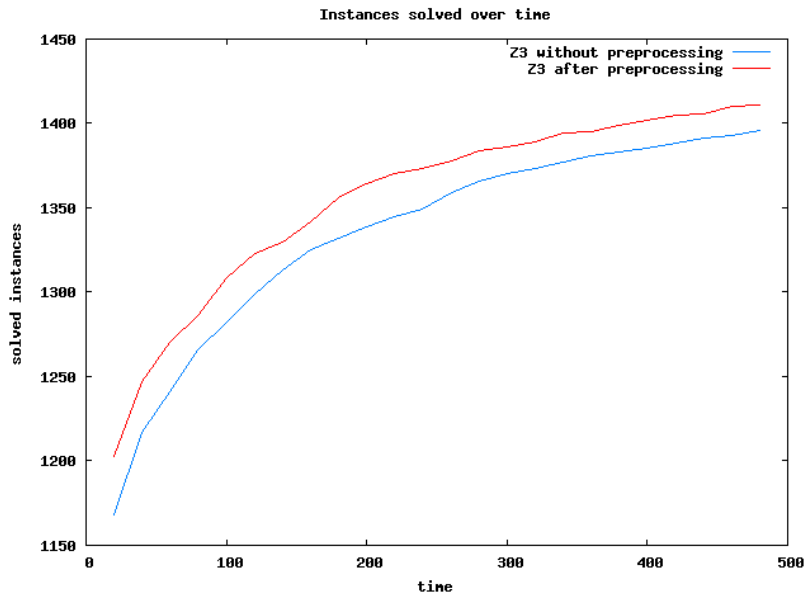
LA simplifications: experiments

Cumulated amount of solved instances: CVC4/smtpp+CVC4



LA simplifications: experiments

Cumulated amount of solved instances: Z3/smtpp+Z3



LA simplifications: user experience

- ▶ The implementer of this module was initially completely unfamiliar with OCaml.
 - ▶ 3-month experience with Scheme
- ▶ One day to get familiar with SMTpp code
 - ▶ core knowledge: AST API
- ▶ A prototype version of the module was implemented in two weeks
 - ▶ quick-and-dirty mode
 - ▶ more effort needed to make it stable

Conclusions

- ▶ SMTpp: already available
 - ▶ supports SMT-LIB 2.5
 - ▶ implemented several features useful to clients of SMT solvers
- ▶ preliminary results cross-checked with other tools
- ▶ usability:
 - ▶ permissive free software license
 - ▶ high-level programming language
- ▶ <https://github.com/ossanha/smtpp>

Future work

- ▶ type checking
- ▶ architectural challenges
 - ▶ processing of very-large benchmarks and memory usage (multiple-passes?)
 - ▶ integration with SMT-solvers (quantifier instantiation)
 - ▶ pluggable architecture: plug-ins *à la* Frama-C
 - ▶ tactic language: user-programmable combination of simplifications?
- ▶ handle requirements such as: traceability, proof production
- ▶ additional simplifications, rewriting:
 - ▶ symmetry detection and symmetry breaking
 - ▶ if-then-else simplifications
- ▶ extend to support delta-debugging, scrambling

Questions?

<https://github.com/ossanha/smtp>

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